

# Computersicherheit (Kryptographie, Authentifizierung, Vertrauen)

Nils Reiter, nils.reiter@uni-koeln.de

October 31, 2023



Was tust Du für Deine Computersicherheit?

bequer

- Was solltest Du tun? Warum tust Du es nicht?
- Welche Aspekte findest Du außerdem relevant?

### Computersicherheit

- Passwörter
- Verschlüsselung (Dateien, E-Mails, ...)
- Sicherheitslücken
- Datensicherheit
- ► Viren/Würmer ...
- Verhalten

...

### Introduction

Not everyone can do/has access to everything

- E.g., only people employed by a university can grade students, and only specific people can see every grade by a specific student
- E.g., only police has access to criminal records (and only in some situations)
- E.g., only fire department can enter and check for emergency exits
- Default case in pre-digital society: Limited access
- Restricted access needs to be modeled digitally as well

### Introduction

Not everyone can do/has access to everything

- E.g., only people employed by a university can grade students, and only specific people can see every grade by a specific student
- E.g., only police has access to criminal records (and only in some situations)
- E.g., only fire department can enter and check for emergency exits
- Default case in pre-digital society: Limited access
- Restricted access needs to be modeled digitally as well
- Multiple aspects
  - Identity and access rights
  - Security of data
  - Trust into systems verifiable behaviour of systems

### Security and Identity

- Closely related technically and conceptually
- If you supply your password, you need to make sure it's to a legitimate party
  - Otherwise, someone could steal it
- No built-in identity verification in many systems (e.g., e-mail)



"On the Internet, nobody knows you're a dog."

Peter Steiner, The New Yorker, 1993

Computersicherheit

- All content can be expressed numerically
  - $\blacktriangleright\,$  E.g., the word "dog" as a  $\langle 100, 111, 103 \rangle$  in ASCII
- ▶ To make things simpler, we only use numbers from now on
- Common example: Alice and Bob want to encrypt stuff, Mallory/Eve are malicious attackers



• Goal: Transform a message m such that only Bob with key k can read it

- Goal: Transform a message m, such that only Bob with key k can read it
- Two main variants
  - Symmetric encryption
  - Asymmetric encryption

### Symmetric Encryption

- ▶ What we do in PDFs, zip-files, ...
- Easiest idea: Key is some number  $\hat{n}$  that we add to each letter to derive at cypher text c

enc(m, n) = m + n  $E.g.: enc(\langle 100, 111, 103 \rangle, 20) = \langle 120, 131, 123 \rangle$   $dec(c, n) \neq c - n$  In this case it's obvious that dec(enc(m, n), n) = m

### Symmetric Encryption

- ▶ What we do in PDFs, zip-files, ...
- $\blacktriangleright$  Easiest idea: Key is some number n that we add to each letter to derive at cypher text c
  - $\blacktriangleright \ \operatorname{enc}(m,n) = m + n$ 
    - E.g.: enc( $\langle \underline{100}, 111, 103 \rangle, 20$ ) =  $\langle \underline{120}, 131, 123 \rangle$
  - $\blacktriangleright \ \operatorname{dec}(c,n) = c \underline{n}$ 
    - ▶ In this case it's obvious that dec(enc(m, n), n) = m
- $\blacktriangleright$  k is a shared secret between Alice and Bob
- They need to agree on k beforehand
  - In a secure way!
- They need to store it and keep it safe
- Anyone with access to k can decrypt the cypher text

### Symmetric Encryption

- Variants
  - More complex calculations (e.g., multiplication instead of addition or complex equations)
    - But they need to be reversible!
  - Encryption of blocks of symbols
  - Generation of more complex keys (e.g., random numbers)
  - Algorithmically adapt n (Enigma)

Enigma simulator

### Symmetric Encryption

- Variants
  - More complex calculations (e.g., multiplication instead of addition or complex equations)
    - But they need to be reversible!
  - Encryption of blocks of symbols
  - Generation of more complex keys (e.g., random numbers)
  - Algorithmically adapt n (Enigma)

### Practical Considerations

- All methods with a constant mapping of symbols are easy to break
- ▶ To regularly change or update *n* requires coordination
- Exchanging a secret is difficult in practice over the internet
- Does not scale well for an entire society

### Encryption Asymmetric Encryption (to produce a shared secret)

- First described to securely generate a shared secret (Diffie/Hellman, 1976; Merkle, 1978)
- "Diffie-Hellman-Key-Exchange"



#### Asymmetric Encryption

**•** Core ingredient: Different keys for encryption  $k_e$  and decryption  $k_d$ 

But generated as a pair, because they need to match

• Requirement:  $\det(enc(m, k_e), k_d) = m$ 

### Asymmetric Encryption

- ▶ Core ingredient: Different keys for encryption  $k_e$  and decryption  $k_e$ 
  - But generated as a pair, because they need to match
- ▶ Requirement:  $dec(enc(m, k_e), k_d) = m$

### **One-Way-Functions**

Functions that are easy to compute in one direction, but impossible to reverse

- Most well known: Prime factorisation
  - Let p and q be two prime numbers
  - Calculating p \* q = i is straightforward and fast
  - But there is no efficient way of getting p and q from i, except iterating over all primes
  - $\blacktriangleright$  If p and q are large enough, this takes until the end of the universe

#### ● ● ● 🗉 🖪 Work ~ 🗧 🔘

W en.wikipedia.org 🔒

.

G

iactorization of HSA-rob, the authors estimate that better algorithms specimeir calculations by a factor of 3-4 and taster computers specimeir calculation by a factor of 1.25-1.67.

#### RSA-250 [edit]

RSA-250 has 250 decimal digits (829 bits), and was factored in February 2020 by Fabrice Boudot, Pierrick Gaudry, Aurore Guillevic, Nadia Heninger, Emmanuel Thomé, and Paul Zimmermann. The announcement of the factorization occurred on February 28.

RSA-250 = 2140324650240744961264423072839333563008614715144755017797754920881418023447 1401366433455190958046796109928518724709145876873962619215573630474547705208 0511905649310668769159001975940569345745223058932597669747168173806936489469 9871578494975937497937

#### RSA-250 = 6413528947707158027879019017057738908482501474294344720811685963202453234463 02386235987526683477087376613925585694639798853367 × 33372027594978156556226010605355114227940760344767555466678452098702384172921 0037080257448673296881877565718986258036032062711

The factorisation of RSA-250 utilised approximately 2700 CPU core-years, using a 2.1 GHz Intel Xeon Gold 6130 CPU as a reference. The computation was performed with the Number Field Sieve algorithm, using the open source CADO-NFS software.

The team dedicated the computation to Peter Montgomery, an American mathematician known for his contributions to computational number theory and cryptography who died on February 18, 2020 and had contributed to factoring RSA-768.<sup>[38]</sup>

#### RSA-260 [edit]



### Encryption Public Key Cryptography

- Core of most modern cryptography
- Public key to encrypt
- Private key to decrypt
- Keys generated as a matching pair
- Used for
  - E-Mails (S/MIME, PGP/GPG)
  - the Web (HTTPS/SSL)
  - SSH



### Encryption Public Key Cryptography

- Public Key Infrastructure (PKI)
- Maps public keys to 'identities'
- For OpenPGP: https://keys.openpgp.org

### Encryption Public Key Cryptography

- Public Key Infrastructure (PKI)
- Maps public keys to 'identities'
- For OpenPGP: https://keys.openpgp.org

	keys.openpgp.org
	We found an entry for nils.reiter@uni-koeln.de.
	https://keys.openpgp.org/vks/v1/by-
	fingerprint/7C210132FD0A0847EFEEE7E6006FF64D646AC7BD
н	int: It's more convenient to use keys.openpgp.org from your OpenPC
	software.
	Take a look at our usage guide for details.

# Signing with Public Key Cryptography

- Public/private keys can also be used for signing messages
- Only Alice is able to produce a message that can be reverted with her public key



# Signing with Public Key Cryptography

- Public/private keys can also be used for signing messages
- Only Alice is able to produce a message that can be reverted with her public key
- Map identities to public keys: "public key infrastructure" (PKI)



# Digital Covid Certificate (EU)

- EU provides central public key infrastructure to sign DCCs
- Member states have their own national systems
  - Public keys of member states are signed on EU level
- Member states can sign public keys of subsidiary elements
  - Theoretically all the way down to individual pharmacies



# Digital Covid Certificate (EU)

- EU provides central public key infrastructure to sign DCCs
- Member states have their own national systems
  - Public keys of member states are signed on EU level
- Member states can sign public keys of subsidiary elements
  - Theoretically all the way down to individual pharmacies



# Digital Covid Certificate (EU)

Distribution of Signing and Validation Information





### Authentication

- How does a system verify a humans identity?
  - Single-factor vs. Multi-factor
- Three general ways all need preparation!
  - Knowledge: Password, PIN, ...
  - Ownership: Hardware token, cell phone, id card, ...
  - ► Inherence: Biometric properties

### **Online-Ausweis**

- Personalausweise enthalten seit 2010 einen RFID-Chip
- Erlaubt maschinelles Auslesen der gespeicherten Daten
- Unterschiede f
  ür hoheitliche und nicht-hoheitliche Funktionen
- Bund signiert Stellen, die z.B. die elD verwenden wollen (PKI etc.)

personalausweisportal.de

### End-To-End-Encryption

- Fully encrypted channel between both proper ends of communication
- Server is not necessarily the 'end'



### End-To-End-Encryption

- Fully encrypted channel between both proper ends of communication
- Server is not necessarily the 'end'
- Communication of people
  - Private devices of communication partners are ends
  - Not the mail server of one of them
- ▶ E2EE ensures that all parties in between the two only see encrypted stuff

### Meta Data

- E-mail content can be encrypted (S/MIME, PGP)
- But meta data is still unencrypted
- That includes



Meta data needed for technical reason, but also leaks information

Questions?

### References I

Diffie, W./M. Hellman (1976). "New directions in cryptography". In: *IEEE Transactions on Information Theory* 22.6, pp. 644–654. DOI: 10.1109/TIT.1976.1055638.
 Merkle, Ralph C. (1978). "Secure Communications over Insecure Channels". In: *Commun. ACM* 21.4, pp. 294–299. ISSN: 0001-0782. DOI: 10.1145/359460.359473. URL: https://doi.org/10.1145/359460.359473.